

## 3.1.1

### HINT 1 (1)

Parity of the bits B1, B2, B3 is computed as:

$$P = B1 \wedge B2 \wedge B3$$

If you don't have a good vectorizing compiler it is faster to use the largest word operation supported by your OS to compute the parity of many bits at a time and to detect wrong words first. Once you have the word you can identify faulty bytes within the word.

Otherwise just do a byte loop !

### **3.1.4**

#### **HINT 1 (1)**

If a byte in stripe 1 is faulty by definition the same byte in all other stripes has to be considered faulty.

Parity alone is not sufficient to detect which stripe is damaged.

## **3.1.4**

### **HINT 1 (2)**

The `sha1sum` program allows you to compute SHA1 sums for files.

## 3.1.5

### HINT 1 (1)

If you have only stripes 1,2,4,P you can recalculate stripe 3 as  
 $3 = 1^2 4^P$

## **3.2.1**

### **HINT 1 (1)**

You might just compare the number of operations needed for  
byte wise XOR or for word XOR operation ....

## **3.2.2**

### **HINT 1 (1)**

Strategy is to recover every row first using diagonal parity, than row-parity aso....

### **3.2.3**

#### **HINT 1 (1)**

If  $X$  is the number of words in a stripe and  $N$  the number of stripes, then  $(N-2)$  is the number of stripes over which the original data gets distributed.

Write down a formula with this symbols or consider the two cases where  $N=4$  and  $N=1.000.000$

## **3.2.4**

### **HINT 1 (1)**

Consider two cases:

- we read the required stripe data and then compute
- we can read and compute in parallel using a sort of read-ahead algorithm for sequential reads